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## Chapter_4 <br> Boolean Algebra (الجبر البوليني)

## Lecture =06

## Booleain Operations \& Laws

(القوانين والععليات الثبولبينةية)

## 6-1. Boolean Operations and Expressions (التعابير والعكليات البولينية)

Variable, complement and literal are terms used in Boolean algebra.
> A variable (عنصر متغير) is a symbol used to represent an action, a condition, or data. A single variable can only have a value of 1 or 0.
$>$ The complement (متهم) represents the inverse of a variable and is indicated with an overbar (خط فوق الرمز).
Thus, the complement of $A$ is $\bar{A}$.
> A literal (بيانات حرفية) is a variable or its complement.

## Boolean Addition operation

Addition is equivalent to the OR operation. The basic rules are shown in Fig.


- In Boolean algebra, a sum term is a sum of literals.
- In logic circuits, a sum term is produced by an OR operation with no AND operations involved. Some examples of sum terms are

$$
A+B ; A+\bar{B} ; A+B+\bar{C} ; \bar{A}+B+C+\bar{D}
$$

## Boolean Multiplication operation

In Boolean algebra, multiplication is equivalent to the AND operation. The basic rules are shown in Fig.


- In Boolean algebra, a product term is the product of literals.
- In logic circuits, a product term is produced by an AND operation with no OR operations involved.
- Some examples of product terms are: $A B ; A \bar{B} ; A B C ; A \bar{B} C \bar{D}$ the expression $\quad A+\bar{B}+C+\bar{D}=0$ ?

For the sum term to be 0 , each literal must $=0$; therefore $A=0, B=1$ (so that $\bar{B}=0$ ), $C=0$, and $D=1$ (so that $\bar{D}=0$ ).

$$
A+\bar{B}+C+\bar{D}=0+\overline{1}+0+\overline{1}=0+0+0+0=0
$$

************************************************************ product term of the expression $A \bar{B} C \bar{D}=1$ ?

SOl|l For the sum term to be 1 , each literal must $=1$; therefore $A=1, B=0$ (so that $\bar{B}=1$ ), $C=1$, and $D=0$ (so that $\bar{D}=1$ ).

$$
A \bar{B} C \bar{D}=1 \cdot \overline{0} \cdot 1 \cdot \overline{0}=1 \cdot 1 \cdot 1 \cdot 1=1
$$

## 6-2. Laws of Boolean Algebra (قو انين الجبر البوليني)

## Commutative Laws (قو انين التبادل)

The commutative laws are applied to addition and multiplication.
$>$ For addition, the commutative law states $\quad A+B=B+A$

$>$ For multiplication, the commutative law states
$A \cdot B=B \cdot A$


## Associative Laws (قو انين التجميع)

$\square$ The associative laws are also applied to addition and multiplication.
$\Rightarrow$ For addition, the commutative law states $\quad A+(B+C)=(A+B)+C$

$>$ For multiplication, the commutative law states $A \cdot(B \cdot C)=(A \cdot B) \cdot C$


## Distributive Law (ققانون التوزيع)

[ The distributive law is the factoring law (قانون العو (قل).
> A common variable can be factored (محالةً إلى عوامل) from an expression just as in ordinary algebra. That is

$$
A \cdot(B+C)=A \cdot B+A \cdot C
$$



## 6-3. Rules of Boolean Algebra (قواعد الجبر البوليني)

Rule_1: $\quad A+0=A$


$$
X=A+0=A
$$

Rule_2: $\quad A+1=1$


$$
X=A+1=1
$$

Rule_3: $\quad$. $\mathbf{0}=\mathbf{0}$

$$
A=1
$$

Rule_4: $\quad A .1=A$
$A=0$


$$
X=A \cdot 1=A
$$

Rules of Boolean Algebra

Rule_5: $\quad A+A=A$


$$
X=A+A=A
$$

Rule_6: $\quad A+\bar{A}=1$


$$
X=A+\bar{A}=1
$$

Rule_7: $\quad \boldsymbol{A} \cdot \boldsymbol{A}=\boldsymbol{A}$


$$
X=A \cdot A=A
$$

Rule_8: $\quad A \cdot \bar{A}=0$

Rule_9: $\quad \overline{\bar{A}}=A$


$$
X=A \cdot \bar{A}=0
$$



$$
\overline{\bar{A}}=A
$$

Rules of Boolean Algebra
Rule_10: $\boldsymbol{A}+\boldsymbol{A B}=\boldsymbol{A}$ This rule can be proved by applying the distributed law, rule_2, and rule_4 as follows:

$$
\begin{aligned}
A+A B & =A(1+B) & & \text { Distributive law } \\
& =A \cdot 1 & & \text { Rul_2: }(1+B)=1 \\
& =A & & \text { Rul_4:A•1=A}
\end{aligned}
$$

The truth table and resulting logic circuit simplification is


Rules of Boolean Algebra
Rule_11: $\boldsymbol{A}+\overline{\boldsymbol{A}} \boldsymbol{B}=\boldsymbol{A}+\boldsymbol{B}$ This rule can be proved as follows:

$$
\begin{aligned}
A+\bar{A} B & =(A+A B)+\bar{A} B & & \text { Rul_10: } A=A+A B \\
& =(A A+A B)+\bar{A} B & & \text { Rul_7: } A=A A \\
& =A A+A B+A \bar{A}+\bar{A} B & & \text { Rul_8:adding } A \bar{A}=0 \\
& =(A+\bar{A})(A+B) & & \text { Distributed law } \\
& =1 \cdot(A+B) & & \text { Rul_6:A+ } \bar{A}=1
\end{aligned}
$$

The truth table and resulting logic circuit simplification is Rule 11: $A+\bar{A} B=A+B$.

| $A$ | $B$ | $\overline{A B}$ | $A+\bar{A} B$ | $A+B$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |  |
| 0 | 1 | 1 | 1 | 1 |  |
| 1 | 0 | 0 | 1 | 1 |  |
| 1 | 1 | 0 | 1 | 1 |  |

Rules of Boolean Algebra
Rule_12: $(A+B)(A+C)=A+B C$

\[

\]

The truth table and resulting logic circuit simplification is

## Rule 12: $(A+B)(A+C)=A+B C$.

| A | B | C | $A+B$ | $A+C$ | $(A+B)(A+C)$ | BC | $A+B C$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 |  | 0 | 0 | 0 | 0 |  |
| 0 | 0 | 1 |  | 1 | 0 | 0 | 0 | ${ }_{B}^{A-0}$ |
| 0 | 1 |  | 1 | 0 | 0 | 0 | 0 |  |
| 0 | 1 |  | 1 | 1 | 1 | 1 | 1 | $C \longrightarrow$ |
| 1 | 0 | 0 | 1 | 1 | 1 | 0 | 1 |  |
| 1 | 0 | 1 | 1 | 1 | 1 | 0 | 1 |  |
| 1 |  | 0 | 1 | 1 | 1 | 0 | 1 | $A$ $B$ $\square$ |
| 1 |  | 1 | 1 | 1 | 1 | 1 | 1 | $C \longrightarrow \square$ |

## Basic rules of Boolean algebra.

1. $A+0=A \quad$ 7. $A \cdot A=A$
2. $A+1=1$
3. $A \cdot \bar{A}=0$
4. $A \cdot 0=0$
5. $\overline{\bar{A}}=A$
6. $A \cdot 1=A$
7. $A+A B=A$
8. $A+A=A$
9. $A+\bar{A} B=A+B$
10. $A+\bar{A}=1 \quad$ 12. $(A+B)(A+C)=A+B C$
$A, B$, or $C$ can represent a single variable or a combination of variables.

## 6-4. DeMorgan's Theorems

## DeMorgan's first Theorem

DeMorgan's first theorem is stated as follows:

- The complement of a product of variables is equal to the sum of the complements of the variables.
- The formula for expressing this theorem for two variables is $\overline{X Y}=\bar{X}+\bar{Y}$



## DeMorgan's second Theorem

DeMorgan's second theorem is stated as follows:

- The complement of a sum of variables is equal to the product of the complements of the variables.
- The formula for expressing this theorem for two variables is $\overline{X+Y}=\bar{X} \cdot \bar{Y}$



# Example 6-3 

Applying DeMorgan's theorems to the expressions

$$
\begin{aligned}
& \overline{X Y Z} ; \quad \overline{X+Y+Z} \\
& \overline{W X Y Z} ; \quad \overline{W+X+Y+Z} \\
& \overline{X Y Z}=\bar{X}+\bar{Y}+\bar{Z} \\
& \overline{W X Y Z}=\bar{W}+\bar{X}+\bar{Y}+\bar{Z} \\
& \overline{X+Y+Z}=\bar{X} \bar{Y} \bar{Z} \\
& \overline{W+X+Y+Z}=\bar{W} \bar{X} \bar{Y} \bar{Z}
\end{aligned}
$$

## Applying DeMorgan's Theorems

The following procedure illustrates the application of DeMorgan's theorems and Boolean algebra to the specific expression

$$
\overline{\overline{A+B \bar{C}}+D(\overline{E+\bar{F}})}
$$

Step 1: Identify the terms to which you can apply DeMorgan's theorems, and think of each term as a single variable. Let $\overline{A+B \bar{C}}=X$ and $D(\overline{E+\bar{F}})=Y$.
Step 2: Since $\overline{X+Y}=\bar{X} \bar{Y}$,

$$
\overline{(\overline{A+B \bar{C}})+(\overline{D(E+\bar{F}})})=(\overline{\overline{A+B} \bar{C}})(\overline{D(\overline{E+\bar{F}})})
$$

Step 3: Use rule $9(\overline{\bar{A}}=A)$ to cancel the double bars over the left term (this is not part of DeMorgan's theorem)

$$
(\overline{\overline{A+B \bar{C}}})(\overline{D(\overline{E+\bar{F}})})=(A+B \bar{C})(\overline{D(\overline{E+\bar{F}})})
$$

Step 4: Apply DeMorgan's theorem to the second term.

$$
(A+B \bar{C})(\overline{D(\overline{E+\bar{F}})})=(A+B \bar{C})(\bar{D}+(\overline{\overline{E+\bar{F}})})
$$

Step 5: Use rule $9(\overline{\bar{A}}=A)$ to cancel the double bars over the $E+\bar{F}$ part of the term.

$$
(A+B \bar{C})(\bar{D}+\overline{\overline{E+\bar{F}}})=(A+B \bar{C})(\bar{D}+E+\bar{F})
$$

## Bample

Apply DeMorgan's theorems to each of the following expressions:
(a) $\overline{(A+B+C) D}$
(b) $\overline{A B C+D E F}$
(c) $A \bar{B}+\bar{C} D+E F$
(a) Let $A+B+C=X$ and $D=Y$. The expression $(\overline{A+B+C) D}$ is of the form $\overline{X Y}=\bar{X}+\bar{Y}$ and can be rewritten as

$$
\overline{(A+B+C) D}=\overline{A+B+C}+\bar{D}
$$

Next, apply DeMorgan's theorem to the term $\overline{A+B+C}$.

$$
\overline{A+B+C}+\bar{D}=\bar{A} \bar{B} \bar{C}+\bar{D}
$$

(b) Let $A B C=X$ and $D E F=Y$ The expression $\overline{A B C+D E F}$ is of the form $\overline{X+Y}=\bar{X} \bar{Y}$ and can be rewritten as

$$
\overline{A B C+D E F}=(\overline{A B C})(\overline{D E F})
$$

Next, apply DeMorgan's theorem to each of the terms $\overline{A B C}$ and $\overline{D E F}$.

$$
(\overline{A B C})(\overline{D E F})=(\bar{A}+\bar{B}+\bar{C})(\bar{D}+\bar{E}+\bar{F})
$$

(c) Let $A \bar{B}=X, \bar{C} D=Y$, and $E F=Z$. The expression $\overline{A \bar{B}+\overline{C D}+E F}$ is of the form $\bar{X}+Y+Z=\bar{X} \bar{Y} \bar{Z}$ and can be rewritten as

$$
\overline{A \bar{B}+\bar{C} D+E F}=(\overline{A \bar{B}})(\overline{\bar{C} D})(\overline{E F})
$$

Next, apply DeMorgan's theorem to each of the terms $\overline{A \bar{B}}, \overline{\bar{C} D}$, and $\overline{E F}$.

$$
(\overline{A \bar{B}})(\overline{\bar{C} D})(\overline{E F})=(\bar{A}+B)(C+\bar{D})(\bar{E}+\bar{F})
$$

## 6-5. Boolean Analysis of Logic Circuits

[] Combinational (التو افقية) logic circuits can be analyzed by writing the expression for each gate and combining the expressions according to the rules for Boolean algebra.
$\square$ For the example, circuit in Fig.


Therefore, the expression for this AND gate is $A(B+C D)$, which is the final output apression for the entire circuit.

## Example 6-5

Given the logic circuit. Apply Boolean algebra to derive the expression for $X$.


## Solution

Write the expression for each gate:

1) For NOR (NOT-OR) gate with $A$ and $B$ inputs we have: $\overline{A+B}$
2) For AND gate with $\overline{A+B}$ and $C$ inputs we have: $C(\overline{A+B})$
3) For OR gate with $C(\overline{A+B})$ and $D$ inputs we have: $C(\overline{A+B})+D$

Therefore, $\quad X=C(\overline{A+B})+D$
Applying DeMorgan's theorems and the distribution law:

$$
\overline{A+B}=\bar{A} \cdot \bar{B} \Rightarrow X=C(\bar{A} \cdot \bar{B})+D=\bar{A} \cdot \bar{B} \cdot C+D
$$

## Constructing a Truth Table for a Logic Circuit



| Inputs |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| $A$ | $B$ | $C$ | $D$ | Output |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |

## 6-6. Logic Simplification Using Boolean Algebra (التّبيط المنطةي)

A simplified Boolean expression uses the fewest gates possible to implement a given expression.
YaII\|\& 6-6 Using Boolean algebra techniques, simplify this expression:

$$
A B+A(B+C)+B(B+C)
$$

Step 1: Apply the distributive law to the second and third terms in the expression, as follows:

$$
A B+A B+A C+B B+B C
$$

Basic rules of Boolean algebra.
Step 2: Apply rule $7(B B=B)$ to the fourth term.

$$
A B+A B+A C+B+B C
$$

Step 3: Apply rule $5(A B+A B=A B)$ to the first two terms.

$$
A B+A C+B+B C
$$

Step 4: Apply rule $10(B+B C=B)$ to the last two terms.

$$
A B+A C+B
$$

| 1. $A+0=A$ | 7. $A \cdot A=A$ |
| :--- | :--- |
| 2. $A+1=1$ | 8. $A \cdot \bar{A}=0$ |
| 3. $A \cdot 0=0$ | 9. $\overline{\bar{A}}=A$ |
| 4. $A \cdot 1=A$ | 10. $A+A B=A$ |
| 5. $A+A=A$ | 11. $A+\bar{A} B=A+B$ |
| 6. $A+\bar{A}=1$ | 12. $(A+B)(A+C)=A+B C$ |

$A, B$, or $C$ can reppesenta single variable or a combination of variables.
Step 5: Apply rule $10(A B+B=B)$ to the first and third terms.

$$
B+A C
$$

$$
A B+A(B+C)+B(B+C) \equiv B+A C
$$

## The simplified circuit



## EMAIIDR S-7 Simplify the following Boolean expression: <br> $$
\bar{A} B C+A \bar{B} \bar{C}+\bar{A} \bar{B} \bar{C}+A \bar{B} C+A B C
$$

Step 1: Factor $B C$ out of the first and last terms.

$$
B C(\bar{A}+A)+A \bar{B} \bar{C}+\bar{A} \bar{B} \bar{C}+A \bar{B} C
$$

Step 2: Apply rule $6(\bar{A}+A=1)$ to the term in parentheses, and factor $A \bar{B}$ from the second and last terms.

$$
B C \cdot 1+A \bar{B}(\bar{C}+C)+\bar{A} \bar{B} \bar{C}
$$

Step 3: Apply rule 4 (drop the 1) to the first term and rule $6(\bar{C}+C=1)$ to the term in parentheses.

$$
B C+A \bar{B} \cdot 1+\bar{A} \bar{B} \bar{C}
$$

Step 4: Apply rule 4 (drop the 1) to the second term.

$$
B C+A \bar{B}+\bar{A} \bar{B} \bar{C}
$$

Step 5: Factor $\bar{B}$ from the second and third terms.

$$
B C+\bar{B}(A+\bar{A} \bar{C})
$$

Step 6: Apply rule $11(A+\bar{A} \bar{C}=A+\bar{C})$ to the term in parentheses.

$$
B C+\bar{B}(A+\bar{C})
$$

Step 7: Use the distributive and commutative laws to get the following expression:

$$
B C+A \bar{B}+\bar{B} \bar{C}
$$

## 6-6. Standard Forms of Boolean Expressions

## (الصيغ النموذجية/ المعيارية للتعابير البولينية)

## The Sum-of-Products (SOP) Form (جمع الجداءات)

$\square$ When two or more product terms are summed by Boolean addition, the resulting expression is a sum-of-products (SOP).
Examples: $\quad A B+A B C ; \quad A B C+C D E+\bar{B} C \bar{D} ; \quad \bar{A} B+\bar{A} B \bar{C}+A C$
AND/OR Implementation of an SOP Expression:

$$
A B+B C D+A C
$$



## NAND/NAND Implementation of an SOP Expression:

- By using only NAND gates, an AND/OR function can be accomplished ( = يُنجز (يتحقق), as illustrated in Figure.
- The first level of NAND gates feed into (يغذي) a NAND gate that acts as a negative-OR gate.

- The NAND and negative-OR inversions cancel and the result is effectively an AND/OR circuit.



## Conversion of General Expression to SOP Form

$\square$ Any logic expression can be changed into SOP form by applying Boolean algebra techniques.
For example, the expression $A(B+C D)$ can be converted to SOP form by applying the distributive law: $A(B+C D)=\mathrm{AB}+\mathrm{ACD}$

Convert each of the following Boolean expressions to SOP form:

$$
A B+B(C D+E F) ; \quad(A+B)(B+C+D) ; \quad \overline{(\overline{A+B})+C}
$$

$$
\begin{aligned}
& A B+B(C D+E F)=A B+B C D+B E F \\
& (A+B)(B+C+D)=A B+A C+A D+B B+B C+B D \\
& \overline{(\overline{A+B})+C}=\overline{(\overline{\overline{A+B}})} \bar{C}=(A+B) \bar{C}=A \bar{C}+B \bar{C}
\end{aligned}
$$

## SOP Standard form

In SOP standard form, every variable in the domain must appear in each term.
Another state is called nonstandard form.
For example:
Standard form: $\quad A \bar{B} C D+\bar{A} \bar{B} C \bar{D}+A B \bar{C} \bar{D}$
Nonstandard form: $\quad A \bar{B} C+\bar{A} \bar{B} \bar{D}+A B \bar{C} \bar{D}$
Converting Product Terms to Standard SOP

- A nonstandard SOP expression can be converted into standard form using Boolean algebra rule $6(A+\bar{A}=1)$.

Step 1: Multiply each nonstandard product term by a term made up of the sum of a missing variable (المتغيرة الصـائعة) and its complement.

Step 2: Repeat Step 1 until all resulting product terms contain all variables in the domain in either complemented or uncomplemented form.

Convert the following Boolean expression into standard SOP

$$
A \bar{B} C+\bar{A} \bar{B}+A B \bar{C} D
$$

The first term, $A \bar{B} C$, is missing variable $D$ or $\bar{D}$, so multiply the first term by $D+\bar{D}$ as follows:

$$
A \bar{B} C=A \bar{B} C(D+\bar{D})=A \bar{B} C D+A \bar{B} C \bar{D}
$$

The second term, $\bar{A} \bar{B}$, is missing variables $C$ or $\bar{C}$ and $D$ or $\bar{D}$, so first multiply the second term by $C+\bar{C}$ as follows: $\bar{A} \bar{B}=\bar{A} \bar{B}(C+\bar{C})=\bar{A} \bar{B} C+\bar{A} \bar{B} \bar{C}$

The two resulting terms are missing variable D or $\bar{D}$, so multiply both terms by $D+\bar{D}$ as follows:

$$
\begin{aligned}
\bar{A} \bar{B} & =\bar{A} \bar{B} C+\bar{A} \bar{B} \bar{C}=\bar{A} \bar{B} C(D+\bar{D})+\bar{A} \bar{B} \bar{C}(D+\bar{D}) \\
& =\bar{A} \bar{B} C D+\bar{A} \bar{B} C \bar{D}+\bar{A} \bar{B} \bar{C} D+\bar{A} \bar{B} \bar{C} \bar{D}
\end{aligned}
$$

The third term, $A B \bar{C} D$, is already in standard form.
The complete standard SOP form of the original expression is as follows:

$$
A \bar{B} C D+A \bar{B} C \bar{D}+\bar{A} \bar{B} C D+\bar{A} \bar{B} C \bar{D}+\bar{A} \bar{B} \bar{C} D+\bar{A} \bar{B} \bar{C} \bar{D}+A B \bar{C} D
$$

## The Product-of-Sums (POS) Form (جداء المجاميح)

$\square$ When two or more sum terms are multiplied, the resulting expression is a product-of-sums (POS).

## Examples

$$
\begin{gathered}
(\bar{A}+B)(A+\bar{B}+C) ; \quad(\bar{A}+\bar{B}+\bar{C})(C+\bar{D}+E)(\bar{B}+C+D) ; \\
(A+B)(A+\bar{B}+C)(\bar{A}+C) ;
\end{gathered}
$$

Implementation of a POS Expression:

$$
(A+B)(B+C+D)(A+C)
$$



## Standard Form of POS

$\square$ In POS standard form, every variable in the domain must appear in each sum term of the expression. Another state is called nonstandard form.
$\square$ For example,
Standard form:

$$
\begin{aligned}
& (\bar{A}+\bar{B}+\bar{C}+\bar{D})(A+\bar{B}+C+D)(A+B+\bar{C}+D) \\
& \quad(\bar{A}+\bar{B}+C)(A+B+\bar{D})(A+\bar{B}+\bar{C}+D)
\end{aligned}
$$

Nonstandard form:
Converting a Sum Term to Standard POS

- A nonstandard POS expression is converted into standard form using Boolean algebra rule $8(A \cdot \bar{A}=0)$.

Step 1: Add to each nonstandard product term a term made up of the product of the missing variable and its complement.

Step 2: Apply rule $12: A+B C=(A+B)(A+C)$.
Step 3: Repeat Step 1 until all resulting sum terms contain all variables in the domain in either complemented or uncomplemented form.

## EAIII\| f-10 Convert the following Boolean expression into standard POS

 form:$$
(A+\bar{B}+C)(\bar{B}+C+\bar{D})(A+\bar{B}+\bar{C}+D)
$$

The first term, $(A+\bar{B}+C)$, is missing variable D or $\bar{D}$, so add $D \bar{D}$ and apply rule 12 as follows: $\quad A+\bar{B}+C+D \bar{D}=(A+\bar{B}+C+D)(A+\bar{B}+C+\bar{D})$
The second term, $(\bar{B}+C+\bar{D})$, is missing variable $A$ or $\bar{A}$, so add $A \bar{A}$ and apply rule 12 as follows: $\quad \bar{B}+C+\bar{D}+A \bar{A}=(A+\bar{B}+C+\bar{D})(\bar{A}+\bar{B}+C+\bar{D})$
The third term, $A+\bar{B}+\bar{C}+D$, is already in standard form.
The standard POS form of the original expression is as follows:

$$
\begin{aligned}
& (A+\bar{B}+C)(\bar{B}+C+\bar{D})(A+\bar{B}+\bar{C}+D)= \\
& (A+\bar{B}+C+D)(A+\bar{B}+C+\bar{D})(A+\bar{B}+C+\bar{D})(\bar{A}+\bar{B}+C+\bar{D})(A+\bar{B}+\bar{C}+D)
\end{aligned}
$$

## Converting SOP Expressions to Truth Table Format

- The first step in constructing a truth table is to list all possible combinations of binary values of the variables in the expression.
- Next, convert the SOP expression to standard form if it is not already.
- Finally, place a 1 in the output column ( $\boldsymbol{X}$ ) for each binary value that makes the standard SOP expression equal to $\mathbf{1}$ and place a $\mathbf{0}$ for all the remaining (المنقفة) binary values.

This procedure is illustrated in Example 6-11.

## Example 6-11

Develop a truth table for the standard SOP expression

$$
\bar{A} \bar{B} C+A \bar{B} \bar{C}+A B C
$$

The binary values that make the product terms in the expressions equal to 1 are

$$
\bar{A} \bar{B} C: 001 ; \quad A \bar{B} \bar{C}: 100 ; \quad A B C: 111
$$

For each of these binary values, place a 1 in the output column as shown in the table. For each of the remaining binary combinations, place a 0 in the output column.

| Inputs |  |  |  | Output |
| :---: | :---: | :---: | :---: | :---: |
| $A$ | $B$ | $C$ | $X$ | Product Term |
| 0 | 0 | 0 | 0 |  |
| 0 | 0 |  | 1 | $\bar{A} \bar{B} C$ |
| 0 |  | 0 | 0 |  |
| 0 | 1 | 1 | 0 |  |
| 1 | 0 | 0 | 1 | $A \bar{B} \bar{C}$ |
| 1 | 0 | 1 | 0 |  |
| 1 | 1 | 0 | 0 |  |
| 1 | 1 | 1 | 1 | $A B C$ |

## Converting POS Expressions to Truth Table Format

To construct a truth table from a POS expression,

- First, list all the possible combinations of binary values of the variables just as was done for the SOP expression.
- Next, convert the POS expression to standard form if it is not already.
- Finally, place a 0 in the output column ( $\boldsymbol{X}$ ) for each binary value that makes the expression equal to 0 and place a 1 for all the remaining binary values. This procedure is illustrated in Example 6-12.

Determine the truth table for the following standard POS

## expression:

$$
(A+B+C)(A+\bar{B}+C)(A+\bar{B}+\bar{C})(\bar{A}+B+\bar{C})(\bar{A}+\bar{B}+C)
$$

The binary values that make the sum terms in the expression equal to 0 are

$$
A+B+C: 000 ; \quad A+\bar{B}+C: 010 ; \quad A+\bar{B}+\bar{C}: 011 ; \bar{A}+B+\bar{C}: 101 ; \bar{A}+\bar{B}+C: 110 ;
$$

For each of these binary values, place a 0 in the output column as shown in the table. For each of the remaining binary combinations, place a 1 in the output column.

| Inputs |  |  | Output |  |
| :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{A}$ | $\boldsymbol{B}$ | $\boldsymbol{C}$ | $\boldsymbol{X}$ | Sum Term |
| 0 | 0 | 0 | 0 | $(A+B+C)$ |
| 0 | 0 | 1 | 1 |  |
| 0 | 1 | 0 | 0 | $(A+\bar{B}+C)$ |
| 0 | 1 | 1 | 0 | $(A+\bar{B}+\bar{C})$ |
| 1 | 0 | 0 | 1 |  |
| 1 | 0 | 1 | 0 | $(\bar{A}+B+\bar{C})$ |
| 1 | 1 | 0 | 0 | $(\bar{A}+\bar{B}+C)$ |
| 1 | 1 | 1 | 1 |  |



## The end of Lecture_06. chapter 4

